

A Construction of Cartesian Closed Categories

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Abstract

In this note we present a simple proof of cartesian closedness of the category of topological spaces and k -continuous maps with the universal property of evaluation maps.

Throughout of this note we assume that \mathcal{C} is a category with finite products, \mathcal{C} is a subcategory of a category \mathcal{A} with $Obj \mathcal{C} = Obj \mathcal{A}$, and the inclusion functor of \mathcal{C} into \mathcal{A} preserves finite products. Moreover, let \mathcal{S} be a full subcategory of \mathcal{C} . We define $\mathcal{S}\mathcal{C}$ as the subcategory of \mathcal{A} with objects all objects of \mathcal{A} and with arrows $f: x \rightarrow y$ all those arrows $f: x \rightarrow y$ in \mathcal{A} for which the composite $f\alpha: s \rightarrow y$ lies in \mathcal{C} for any arrow $\alpha: s \rightarrow x$ in \mathcal{C} with $s \in Obj \mathcal{S}$. We can easily verify that \mathcal{C} is a subcategory of $\mathcal{S}\mathcal{C}$ and the inclusion functor of \mathcal{C} into $\mathcal{S}\mathcal{C}$ preserves finite products. We say an object x in \mathcal{A} is an \mathcal{S} -generated object if any arrow $f: x \rightarrow y$ in $\mathcal{S}\mathcal{C}$ with domain x lies in \mathcal{C} . Let $\mathcal{S}\mathcal{C}$ denote the full subcategory of \mathcal{C} with objects all \mathcal{S} -generated objects.

Let $k: \mathcal{S}\mathcal{C}^{op} \times \mathcal{S}\mathcal{C} \rightarrow \mathcal{A}$ be a functor with a dinatural transformation $\varepsilon_{\langle y, z \rangle}: k(y, z) \times y \rightarrow z$ in \mathcal{A} (precisely, natural in z and dinatural in y), that is, for any arrows $g: z \rightarrow z'$ and $h: y' \rightarrow y$ in $\mathcal{S}\mathcal{C}$, the following two diagrams

$$\begin{array}{ccc}
 & k(y, z) \times y & \xrightarrow{\varepsilon_{\langle y, z \rangle}} z \\
 (1) \quad & \downarrow k(y, g) \times y & \downarrow g \\
 & k(y, z') \times y & \xrightarrow{\varepsilon_{\langle y, z' \rangle}} z'
 \end{array}$$

$$\begin{array}{ccc}
 & k(y, z) \times y' & \xrightarrow{k(y, z) \times h} k(y, z) \times y \\
 (2) \quad & \downarrow k(h, z) \times y' & \downarrow \varepsilon_{\langle y, z \rangle} \\
 & k(y', z) \times y' & \xrightarrow{\varepsilon_{\langle y', z \rangle}} z
 \end{array}$$

are commutative. (Note that $a \times b$ denotes the product of a and b in \mathcal{C} .)

The dinatural transformation $\varepsilon_{\langle y, z \rangle}: k(y, z) \times y \rightarrow z$ is called *quasi-universal* if it satisfies the following universal property: For any arrow $f: x \times y \rightarrow z$ in $\mathcal{S}\mathcal{C}$, there is a unique arrow $\hat{f}: x \rightarrow k(y, z)$ in \mathcal{A} such that $f = \varepsilon_{\langle y, z \rangle}(\hat{f} \times y)$.

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$$(3) \quad \begin{array}{ccc} x & & x \times y \\ \downarrow f & & \downarrow f \times y \quad \searrow f \\ k(y, z) & & k(y, z) \times y \xrightarrow{\varepsilon_{\langle y, z \rangle}} z \end{array}$$

An object y in \mathcal{A} is (\mathcal{S} -) *admissible* if $\varepsilon_{\langle y, z \rangle}$ is an arrow in $(\mathcal{S})\mathcal{C}$ for any object z in \mathcal{A} . An object y in \mathcal{A} is (\mathcal{S} -) *proper* if, for any arrow $f: x \times y \rightarrow z$ in $(\mathcal{S})\mathcal{C}$, the unique arrow $\hat{f}: x \rightarrow k(y, z)$ is an arrow in $(\mathcal{S})\mathcal{C}$, that is, if, whenever $\varepsilon_{\langle y, z \rangle}(h \times y)$ lies in $(\mathcal{S})\mathcal{C}$, so does h .

Throughout the rest of this note we assume that there are given a functor $k: \mathcal{S}\mathcal{C}^{op} \times \mathcal{S}\mathcal{C} \rightarrow \mathcal{A}$ with a quasi-universal dinatural transformation $\varepsilon_{\langle y, z \rangle}: k(y, z) \times y \rightarrow z$ and a full subcategory \mathcal{U} of \mathcal{C} with $\mathcal{S} \subset \mathcal{U} \subset \mathcal{S}\mathcal{C}$, satisfying the following four axioms:

- (A) If s and t are two objects in \mathcal{S} , then $s \times t$ is an \mathcal{S} -generated object.
- (B) Every object s in \mathcal{S} is admissible.
- (C) Every object u in \mathcal{U} is proper.
- (D) An arrow $h: x \rightarrow k(y, z)$ in \mathcal{A} is an arrow in \mathcal{C} if and only if, for any arrow $n: u \rightarrow y$ in \mathcal{C} with $u \in \text{Obj}\mathcal{U}$, the composite

$$x \xrightarrow{h} k(y, z) \xrightarrow{k(n, z)} k(u, z)$$

lies in \mathcal{C} .

Lemma 1. For any arrow $n: u \rightarrow y$ in \mathcal{C} with $u \in \text{Obj}\mathcal{U}$, the arrow $k(n, z): k(y, z) \rightarrow k(u, z)$ lies in \mathcal{C} .

Proof. It is immediate from the axiom (D).

Proposition 2. If $g: y' \rightarrow y$ is an arrow in $\mathcal{S}\mathcal{C}$, then $k(g, z): k(y, z) \rightarrow k(y', z)$ is an arrow in \mathcal{C} .

Proof. By virtue of the axiom (D), we have only to prove that the composite $k(n, z)k(g, z)$ lies in \mathcal{C} for any arrow $n: u \rightarrow y'$ in \mathcal{C} with $u \in \text{Obj}\mathcal{U}$. But $k(n, z)k(g, z) = k(gn, z)$, that is, the triangle

$$\begin{array}{ccc} k(y, z) & \xrightarrow{k(g, z)} & k(y', z) \\ & \searrow k(gn, z) & \downarrow k(n, z) \\ & & k(u, z) \end{array}$$

is commutative and gn lies in \mathcal{C} , because u is an \mathcal{S} -generated object. Hence, by Lemma 1, $k(gn, z)$ lies in \mathcal{C} .

Theorem 3. Every object y in \mathcal{A} is \mathcal{S} -admissible.

Proof. We have to show that $\varepsilon_{\langle y, z \rangle} \langle \alpha, \beta \rangle$ lies in \mathcal{C} for any arrow $\langle \alpha, \beta \rangle: s \rightarrow k(y, z) \times y$ in \mathcal{C} with $s \in \text{Obj}\mathcal{S}$. In the commutative diagram

$$\begin{array}{ccccc}
 s & \xrightarrow{\langle \alpha, \beta \rangle} & k(y, z) \times y & \xrightarrow{\varepsilon_{\langle y, z \rangle}} & z \\
 & \searrow \langle \alpha, s \rangle & \uparrow k(y, z) \times \beta & & \uparrow \varepsilon_{\langle y, z \rangle} \\
 & & k(y, z) \times s & \xrightarrow{k(\beta, z) \times s} & k(s, z) \times s,
 \end{array}$$

three arrows $\langle \alpha, s \rangle$, $k(\beta, z) \times s$ and $\varepsilon_{\langle s, z \rangle}$ lie in \mathcal{C} , by Lemma 1 ($s \in \mathcal{S} \subset \mathcal{U}$) and the axiom (B). Therefore $\varepsilon_{\langle y, z \rangle} \langle \alpha, \beta \rangle$ is an arrow in \mathcal{C} , as desired.

Theorem 4. *If u is an \mathcal{S} -generated object and s is an object in \mathcal{S} , then $u \times s$ is an \mathcal{S} -generated object.*

Proof. To prove this theorem, it suffices to show that every arrow $f: u \times s \rightarrow z$ in $\mathcal{S}\mathcal{C}$ lies in \mathcal{C} . Let $\alpha: t \rightarrow u$ be an arrow in \mathcal{C} with $t \in \text{Obj } \mathcal{S}$. In the commutative diagram in \mathcal{A}

$$\begin{array}{ccc}
 t \times s & \xrightarrow{\alpha \times s} & u \times s \\
 \alpha \hat{f} \times s \searrow & & \downarrow f \times s \\
 & & k(s, z) \times s \xrightarrow{\varepsilon_{\langle s, z \rangle}} z,
 \end{array}$$

$f(\alpha \times s)$ lies in $\mathcal{S}\mathcal{C}$ ($\alpha \times s \in \mathcal{C}$, $f \in \mathcal{S}\mathcal{C}$) and hence $f(\alpha \times s)$ lies in \mathcal{C} since $t \times s$ is an \mathcal{S} -generated object from the axiom (A). Applying the axiom (C) to $s \in \mathcal{S} \subset \mathcal{U}$, we have $\alpha \hat{f}$ is an arrow in \mathcal{C} . This states that \hat{f} lies in $\mathcal{S}\mathcal{C}$. But u is an \mathcal{S} -generated object, so \hat{f} lies in \mathcal{C} . From the axiom (B), $\varepsilon_{\langle s, z \rangle}$ lies in \mathcal{C} and consequently so does $f = \varepsilon_{\langle s, z \rangle}(\hat{f} \times s)$. This proves the theorem.

Theorem 5. *Every object y in \mathcal{A} is \mathcal{S} -proper.*

Proof. To prove this theorem we will show that, for any arrow $f: x \times y \rightarrow z$ in $\mathcal{S}\mathcal{C}$, the unique arrow $\hat{f}: x \rightarrow k(y, z)$ lies in $\mathcal{S}\mathcal{C}$. Consider the commutative diagram

$$\begin{array}{ccccc}
 s \times u & \xrightarrow{\alpha \times n} & & & x \times y \\
 & \searrow f \alpha \times u & & & \searrow f \times y \\
 & & k(y, z) \times u & \xrightarrow{k(y, z) \times n} & k(y, z) \times y \\
 k(n, z) \hat{f} \alpha \times u \downarrow & & \swarrow k(n, z) \times u & & \swarrow \varepsilon_{\langle y, z \rangle} \\
 k(u, z) \times u & \xrightarrow{\varepsilon_{\langle u, z \rangle}} & & & z
 \end{array}$$

for all arrows $\alpha: s \rightarrow x$ in \mathcal{C} with $s \in \text{Obj } \mathcal{S}$ and $n: u \rightarrow y$ in \mathcal{C} with $u \in \text{Obj } \mathcal{U}$. Since $s \times u$ is an \mathcal{S} -generated object by Theorem 4, $f(\alpha \times n)$ lies in \mathcal{C} . On the other hand, u is proper from the axiom (C). Hence $k(n, z) \hat{f} \alpha$ lies in \mathcal{C} and $\hat{f} \alpha$ lies in \mathcal{C} from the axiom (D). This proves that \hat{f} lies in $\mathcal{S}\mathcal{C}$.

Theorem 6. *If $g: z \rightarrow z'$ is an arrow in $\mathcal{S}\mathcal{C}$, then $k(y, g): k(y, z) \rightarrow k(y, z')$ is an arrow in $\mathcal{S}\mathcal{C}$.*

Proof. Consider the commutative square

$$\begin{array}{ccc}
 k(y, z) \times y & \xrightarrow{\varepsilon_{\langle y, z \rangle}} & z \\
 k(y, g) \times y \downarrow & & \downarrow g \\
 k(y, z') \times y & \xrightarrow{\varepsilon_{\langle y, z' \rangle}} & z'.
 \end{array}$$

Since g is an arrow in $\mathcal{S}\mathcal{C}$ from the hypothesis and $\varepsilon_{\langle y, z \rangle}$ is an arrow in $\mathcal{S}\mathcal{C}$ by Theorem 3, the composite $g\varepsilon_{\langle y, z \rangle}$ lies in $\mathcal{S}\mathcal{C}$ and so the arrow $k(y, g) = \{g\varepsilon_{\langle y, z \rangle}\}^\wedge$ lies in $\mathcal{S}\mathcal{C}$ by Theorem 5, as desired.

As a consequence of Proposition 2, Theorem 6 and Theorem 3, we can conclude the following

Theorem 7. *The category $\mathcal{S}\mathcal{C}$ is cartesian closed.*

From the standard arguments in cartesian closed categories, we have the exponential laws: For all objects x, y, z in \mathcal{C} , there exist natural isos

$$k(x \times y, z) \simeq k(x, k(y, z))$$

and

$$k(x, y \times z) \simeq k(x, y) \times k(x, z)$$

in $\mathcal{S}\mathcal{C}$.

Example 8. Let \mathcal{C} be the category of topological spaces and continuous maps, \mathcal{A} the category of topological spaces and set maps, and \mathcal{S} the full subcategory of \mathcal{C} consisting of compact Hausdorff spaces. Then $\mathcal{S}\mathcal{C}$ is the category of compactly generated spaces [5]. Further, let \mathcal{U} be some full subcategory of \mathcal{C} with $\mathcal{S} \subset \mathcal{U} \subset \mathcal{S}\mathcal{C}$. Denote by $F(X, Y)$ the function space of all continuous maps $X \rightarrow Y$ with compact open topology, and by $k(Y, Z)$ the function space of all maps $Y \rightarrow Z$ in $\mathcal{S}\mathcal{C}$ with the initial (or smallest) topology determined by the family of set maps

$$k(Y, Z) \xrightarrow{n^*} F(U, Z): f \mapsto fn,$$

where $U \in \text{Obj } \mathcal{U}$ and $n: U \rightarrow Y$ is a continuous map. (Note that n^* is well defined because $U \in \text{Obj } \mathcal{S}\mathcal{C}$.) It is trivial that $k(U, Z) = F(U, Z)$ for $U \in \text{Obj } \mathcal{U}$, $k: \mathcal{S}\mathcal{C}^{\text{op}} \times \mathcal{S}\mathcal{C} \rightarrow \mathcal{A}$ is a bifunctor and the axiom (D) is satisfied. Next, define $\varepsilon_{\langle Y, Z \rangle}: k(Y, Z) \times Y \rightarrow Z$ to be the usual evaluation map. Then it can be checked without difficulty that $\varepsilon_{\langle Y, Z \rangle}$ is a quasi-universal dinatural transformation and that the axioms (A), (B) and (C) follows from the familiar properties of compact open topology. Hence, the category $\mathcal{S}\mathcal{C}$ is a cartesian closed category by Theorem 7.

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